## AMENDMENTS TO THE CLAIMS

1. (Original) A method for assigning a channel to a UE (user equipment) by a UTRAN (UMTS (Universal Mobile Telecommunications System) Terrestrial Radio Access Network) in a CDMA (Code Division Multiple Access) communication system, the method comprising the steps of:

receiving a access preamble signature from the UE; and

selecting one of a plurality of channel assignment signatures associated with the received access preamble signature in order to assign one of a plurality of physical common packet channels (PCPCHs) unused in the UTRAN.

- 2. (Original) The method as claimed in claim 1, wherein the UTRAN selects one of the channel assignment signatures depending on a maximum data rate required when the UE transmits data.
- 3. (Currently Amended) The method as claimed in claim 1, further comprising the step of selecting one of the PCPCHs unused in the UTRAN depending on the received access preamble signature and the selected channel assignment signature for receiving a packet data from the UE.
- 4. (Original) The method as claimed in claim 3, wherein the PCPCH selecting step comprises the steps of:

determining a number P<sub>SF</sub> of PCPCHs capable of supporting a maximum data rate required when the UE transmits data out of the unused PCPCHs;

determining a number  $S_{SF}$  of access preamble signature available for the maximum data rate required when the UE transmits data;

determining a number  $T_{SF}$  of channel assignment signatures available for the maximum data rate depending on the number  $P_{SF}$  of the PCPCHs;

calculating a minimum positive number  $M_{SF}$  out of positive numbers which are determined to have a remainder of '0' when multiplying the number  $S_{SF}$  of the access preamble signatures by a given positive number and dividing the multiplied value by the number  $P_{SF}$  of the PCPCHs;

calculating a specific coefficient 'n' satisfying the following equation

$$n*M_{SF}*S_{SF} \le i+j*S_{SF} < (n+1)*M_{SF}*S_{SF}$$

where i denotes an access preamble signature number and i denotes a channel allocation message number; and

selecting one PCPCH's number 'k' out of the PCPCHs unused in the UTRAN by satisfying the following equation

$$k = \{[(i+n) \mod S_{SF}] + j*S_{SF}\} \mod P_{SF}.$$

(Original) The method as claimed in claim 4, further comprising the steps of: 5. calculating a specific coefficient 'm' for determining a data rate by satisfying the following equation

$$P_{2^{m-1}} \leq k < P_{2^m}$$

where  $P_{2^{m-1}}$  denotes a channelization code with a spreading factor  $2^{m-1}$ , and  $P_{2^m}$  denotes a channelization code with a spreading factor 2<sup>m</sup>;

calculating an uplink scrambling code's number by satisfying the following equation

$$\left[ \sum_{2 \le a < m-1} (P_{2^a} - P_{2^{a-1}}) / 2^{a-1} + (k - P_{2^{m-1}}) / 2^m \right]$$

where,  $\alpha$  is an integer numbers

calculating a heading node by satisfying the following equation 
$$\left(\sum_{2\leq a\leq m-1}(P_{2^a}-P_{2^{a-1}})*2^{m-a}+k-P_{2^{m-1}}\right)/2^{m-1}$$

; and

selecting a channelization code with a spreading factor corresponding to the maximum data rate from the heading node and determining the selected channelization code as a channelization code to be used by the UE.

6. (Original) The method as claimed in claim 1, wherein the channel assignment signature (i) is selected by satisfying following equation;

$$n*M_{SF}*S_{SF} \le i+j*S_{SF} < (n+1)*M_{SF}*S_{SF}$$

where, i is number of the access preamble signature, the S<sub>SF</sub> is a number of access preamble signatures assigned for the maximum data rate determined by the access preamble signature, the M<sub>SF</sub> is a minimum positive number(MSF) out of positive numbers which are determined to have a remainder of '0' when multiplying the number S<sub>SF</sub> by a given positive number and dividing the multiplied value by a number  $P_{SF}$  representing number of PCPCHs assigned to support the maximum data rate, the n indicates how many times a period of  $M_{SF}$  has been repeated.

7. (Original) The method as claimed in claim 6, wherein a PCPCH (k) is determined by satisfying following equation;

$$k = \{[(i+n) \text{ mod } S_{SF}] + j*S_{SF}\} \text{ mod } P_{SF}.$$

8. (Original) A method for assigning a channel to a UE (user equipment) by a UTRAN (UMTS (Universal Mobile Telecommunications System) Terrestrial Radio Access Network) in a CDMA (Code Division Multiple Access) communication system, the method comprising the steps of:

receiving a selected one of a plurality of access preamble signatures from the UE; and determining a specific channel assignment signature from a plurality of channel assignment signatures so as to select one of a plurality of unused PCPCHs (physical common packet channels) depending on the received access preamble signature and a channel assignment signature.

- 9. (Original) The method as claimed in claim 8, wherein the UTRAN selects one of the channel assignment signatures depending on a maximum data rate determined by the access preamble signature.
- 10. (Original) The method as claimed in claim 9, wherein the channel assignment signature (j) is selected by satisfying following equation;

$$n^*M_{SF}^*S_{SF} \leq i + j^*S_{SF} < (n+1)^*M_{SF}^*S_{SF}$$

where, i is number of the access preamble signature, the  $S_{SF}$  is a number of access preamble signatures assigned for the maximum data rate determined by the access preamble signature, the  $M_{SF}$  is a minimum positive number  $(M_{SF})$  out of positive numbers which are determined to have a remainder of '0' when multiplying the number  $S_{SF}$  by a given positive number and dividing the multiplied value by a number  $P_{SF}$  representing number of PCPCHs assigned to support the maximum data rate and the n indicates how many times a period of  $M_{SF}$  has been repeated.

- (Original) The method as claimed in claim 10, further comprising the step of selecting one of the PCPCHs unused in the UTRAN depending on the received access preamble signature and the selected channel assignment signature for receiving a packet data from the UE.
- 12. (Currently Amended) The method as claimed in claim 11, wherein <u>a</u> the selected PCPCH (k) is determined by satisfying following equation;

$$k = \{[(i+n) \mod S_{SF}] + j*S_{SF}\} \mod P_{SF}.$$

13. (Original) The method as claimed in claim 9, wherein the PCPCH selecting step comprises the steps of:

determining a number P<sub>SF</sub> of PCPCHs capable of supporting a maximum data rate required when the UE transmits data out of the unused PCPCHs;

determining a number  $S_{SF}$  of access preamble signatures available for the maximum data rate required when the UE transmits data;

determining a number  $T_{SF}$  of channel assignment signatures available for the maximum data rate depending on the number  $P_{SF}$  of the PCPCHs;

calculating a minimum positive number  $M_{SF}$  out of positive numbers which are determined to have a remainder of '0' when multiplying the number  $S_{SF}$  of the access preamble signatures by a given positive number and dividing the multiplied value by the number  $P_{SF}$  of the PCPCHs;

calculating a specific coefficient 'n' satisfying the following equation

$$n*M_{SF}*S_{SF} \le i+j*S_{SF} < (n+1)*M_{SF}*S_{SF}$$

where i denotes an access preamble signature number and j denotes a channel allocation message number; and

selecting one PCPCH's number 'k' out of the PCPCHs unused in the UTRAN by satisfying the following equation

$$k = \{[(i+n) \mod S_{SF}] + j*S_{SF}\} \mod P_{SF}.$$

14. (Original) The method as claimed in claim 13, further comprising the steps of: calculating a specific coefficient 'm' for determining a data rate by satisfying the following equation

$$P_{2^{m-1}} \leq k < P_{2^m}$$

where  $P_{2^{m-1}}$  denotes a channelization code with a spreading factor  $2^{m-1}$ , and  $P_{2^m}$  denotes a channelization code with a spreading factor 2<sup>m</sup>;

calculating an uplink scrambling code's number by satisfying the following equation

calculating an uplink scrambling code's number by satisfying
$$\left[\sum_{2 \le a < m-1} (P_{2^a} - P_{2^{a-1}})/2^{a-1} + (k - P_{2^{m-1}})/2^m\right]$$
where, a is an integer numbers;

calculating a heading node by satisfying the following equation

$$\left(\sum_{2 \le a \le m-1} (P_{2^a} - P_{2^{a-1}}) * 2^{m-a} + k - P_{2^{m-1}}\right) / 2^{m-1}$$

; and

selecting a channelization code with a spreading factor corresponding to the maximum data rate from the heading node and determining the selected channelization code as a channelization code to be used by the UE.

(Original) A method for assigning a channel in a UE (user equipment) for a CDMA 15. (Code Division Multiple Access) communication system, comprising the steps of:

upon generation of data to be transmitted over a PCPCH channel, selecting one of a plurality of access preamble signatures and transmitting the selected access preamble signature to a UTRAN;

receiving a selected one of a plurality of channel assignment signatures from the UTRAN; and

determining a PCPCH channel for transmitting the data depending on the selected access preamble signature and the received channel assignment signature.

- 16. (Original) The method as claimed in claim 15, wherein the UE selects one of the access preamble signatures depending on a maximum data rate required when transmitting the data.
- (Currently Amended) The method as claimed in claim 15, wherein the PCPCH (k) is 17. determined by satisfying following equation;

$$k = \{[(i+n) \text{ mod } S_{SF}] + j*S_{SF}\} \text{ mod } P_{SF}[[.]]$$

where, i is a number of the access preamble signature, the j is a number of the received channel assignment signature, the S<sub>SF</sub> is a number of access preamble signatures assigned for the maximum data rate determined by the access preamble signature, the P<sub>SF</sub> representing a number of PCPCHs assigned to support the maximum data rate, and the n indicates how many times a period of  $M_{SF}$ , which represent a minimum positive number out of positive numbers which are determined to have a remainder of '0' when multiplying the number  $S_{SF}$  by a given positive number and dividing the multiplied value by a number  $P_{SF}$ , has been repeated.

18. (Original) The method as claimed in claim 15, wherein the selecting step comprises the steps of:

determining a number P<sub>SF</sub> of PCPCHs capable of supporting a maximum data rate required when the UE transmits data out of the unused PCPCHs;

determining a number  $S_{SF}$  of access preamble signatures available for the maximum data rate required when the UE transmits data;

determining a number  $T_{SF}$  of channel assignment signatures available for the maximum data rate depending on the number  $P_{SF}$  of the PCPCHs;

calculating a minimum positive number  $M_{SF}$  out of positive numbers which are determined to have a remainder of '0' when multiplying the number  $S_{SF}$  of the access preamble signatures by a given positive number and dividing the multiplied value by the number  $P_{SF}$  of the PCPCHs;

calculating a specific coefficient 'n' satisfying the following equation

$$n*M_{SF}*S_{SF} \le i+j*S_{SF} < (n+1)*M_{SF}*S_{SF}$$

where i denotes an access preamble signature number and j denotes a channel allocation message number; and

selecting one PCPCH's number 'k' out of the PCPCHs unused in the UTRAN by satisfying the following equation

$$k = \{[(i+n) \mod S_{SF}] + j*S_{SF}\} \mod P_{SF}.$$

19. (Original) The method as claimed in claim 18, further comprising the steps of: calculating a specific coefficient 'm' for determining a data rate by satisfying the following equation

$$P_{2^{m-1}} \leq k < P_{2^m}$$

where  $P_{2^{m-1}}$  denotes a channelization code with a spreading factor  $2^{m-1}$ , and  $P_{2^m}$  denotes a channelization code with a spreading factor  $2^m$ ;

calculating an uplink scrambling code's number by satisfying the following equation

$$\left[ \sum_{2 \le a < m-1} (P_{2^a} - P_{2^{a-1}}) / 2^{a-1} + (k - P_{2^{m-1}}) / 2^m \right]$$

where, a is an integer numbers;

calculating a heading node by satisfying the following equation 
$$\left(\sum_{2\leq a\leq m-1} (P_{2^a}-P_{2^{a-1}})*2^{m-a}+k-P_{2^{m-1}}\right)/2^{m-1}$$

; and

selecting a channelization code with a spreading factor corresponding to the maximum data rate from the heading node and determining the selected channelization code as a channelization code to be used by the UE.